
Linux Hid Documentation

The kernel development community

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CONTENTS

CARE AND FEEDING OF YOUR HUMAN INTERFACE DEVICES

1.1 Introduction

In addition to the normal input type HID devices, USB also uses the human interface device protocols for things that are not really human interfaces, but have similar sorts of communication needs. The two big examples for this are power devices (especially uninterruptable power supplies) and monitor control on higher end monitors.

To support these disparate requirements, the Linux USB system provides HID events to two separate interfaces: * the input subsystem, which converts HID events into normal input device interfaces (such as keyboard, mouse and joystick) and a normalised event interface - see Documentation/input/input.rst * the hiddev interface, which provides fairly raw HID events

The data flow for a HID event produced by a device is something like the following:

```
usb.c ---> hid-core.c ----> hid-input.c ----> [keyboard/mouse/joystick/  
->event]  
|  
--> hiddev.c ----> POWER / MONITOR CONTROL
```

In addition, other subsystems (apart from USB) can potentially feed events into the input subsystem, but these have no effect on the hid device interface.

1.2 Using the HID Device Interface

The hiddev interface is a char interface using the normal USB major, with the minor numbers starting at 96 and finishing at 111. Therefore, you need the following commands:

```
mknod /dev/usb/hiddev0 c 180 96  
mknod /dev/usb/hiddev1 c 180 97  
mknod /dev/usb/hiddev2 c 180 98  
mknod /dev/usb/hiddev3 c 180 99  
mknod /dev/usb/hiddev4 c 180 100  
mknod /dev/usb/hiddev5 c 180 101  
mknod /dev/usb/hiddev6 c 180 102
```

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```
mknod /dev/usb/hiddev7 c 180 103
mknod /dev/usb/hiddev8 c 180 104
mknod /dev/usb/hiddev9 c 180 105
mknod /dev/usb/hiddev10 c 180 106
mknod /dev/usb/hiddev11 c 180 107
mknod /dev/usb/hiddev12 c 180 108
mknod /dev/usb/hiddev13 c 180 109
mknod /dev/usb/hiddev14 c 180 110
mknod /dev/usb/hiddev15 c 180 111
```

So you point your hiddev compliant user-space program at the correct interface for your device, and it all just works.

Assuming that you have a hiddev compliant user-space program, of course. If you need to write one, read on.

1.3 The HIDDEV API

This description should be read in conjunction with the HID specification, freely available from <http://www.usb.org>, and conveniently linked of <http://www.linux-usb.org>.

The hiddev API uses a `read()` interface, and a set of `ioctl()` calls.

HID devices exchange data with the host computer using data bundles called “reports”. Each report is divided into “fields”, each of which can have one or more “usages”. In the hid-core, each one of these usages has a single signed 32 bit value.

1.3.1 `read()`:

This is the event interface. When the HID device’s state changes, it performs an interrupt transfer containing a report which contains the changed value. The `hid-core.c` module parses the report, and returns to `hiddev.c` the individual usages that have changed within the report. In its basic mode, the hiddev will make these individual usage changes available to the reader using a struct `hiddev_event`:

```
struct hiddev_event {
    unsigned hid;
    signed int value;
};
```

containing the HID usage identifier for the status that changed, and the value that it was changed to. Note that the structure is defined within `<linux/hiddev.h>`, along with some other useful `#defines` and structures. The HID usage identifier is a composite of the HID usage page shifted to the 16 high order bits ORed with the usage code. The behavior of the `read()` function can be modified using the `HIDIOCSFLAG` `ioctl()` described below.

1.3.2 ioctl():

This is the control interface. There are a number of controls:

HIDIOCGVERSION

- int (read)

Gets the version code out of the hiddev driver.

HIDIOCAPPLICATION

- (none)

This ioctl call returns the HID application usage associated with the hid device. The third argument to ioctl() specifies which application index to get. This is useful when the device has more than one application collection. If the index is invalid (greater or equal to the number of application collections this device has) the ioctl returns -1. You can find out beforehand how many application collections the device has from the num_applications field from the hiddev_devinfo structure.

HIDIOCGCOLLECTIONINFO

- struct hiddev_collection_info (read/write)

This returns a superset of the information above, providing not only application collections, but all the collections the device has. It also returns the level the collection lives in the hierarchy. The user passes in a hiddev_collection_info struct with the index field set to the index that should be returned. The ioctl fills in the other fields. If the index is larger than the last collection index, the ioctl returns -1 and sets errno to -EINVAL.

HIDIOCGDEVINFO

- struct hiddev_devinfo (read)

Gets a hiddev_devinfo structure which describes the device.

HIDIOCGSTRING

- struct hiddev_string_descriptor (read/write)

Gets a string descriptor from the device. The caller must fill in the “index” field to indicate which descriptor should be returned.

HIDIOCINITREPORT

- (none)

Instructs the kernel to retrieve all input and feature report values from the device. At this point, all the usage structures will contain current values for the device, and will maintain it as the device changes. Note that the use of this ioctl is unnecessary in general, since later kernels automatically initialize the reports from the device at attach time.

HIDIOCGNAME

- string (variable length)

Gets the device name

HIDIOCGREPORT

- struct `hiddev_report_info` (write)

Instructs the kernel to get a feature or input report from the device, in order to selectively update the usage structures (in contrast to `INITREPORT`).

HIDIOCSREPORT

- struct `hiddev_report_info` (write)

Instructs the kernel to send a report to the device. This report can be filled in by the user through `HIDIOCSUSAGE` calls (below) to fill in individual usage values in the report before sending the report in full to the device.

HIDIOCGREPORTINFO

- struct `hiddev_report_info` (read/write)

Fills in a `hiddev_report_info` structure for the user. The report is looked up by type (input, output or feature) and id, so these fields must be filled in by the user. The ID can be absolute - the actual report id as reported by the device - or relative - `HID_REPORT_ID_FIRST` for the first report, and `(HID_REPORT_ID_NEXT | report_id)` for the next report after `report_id`. Without a-priori information about report ids, the right way to use this ioctl is to use the relative IDs above to enumerate the valid IDs. The ioctl returns non-zero when there is no more next ID. The real report ID is filled into the returned `hiddev_report_info` structure.

HIDIOCGFIELDINFO

- struct `hiddev_field_info` (read/write)

Returns the field information associated with a report in a `hiddev_field_info` structure. The user must fill in `report_id` and `report_type` in this structure, as above. The `field_index` should also be filled in, which should be a number from 0 and `maxfield-1`, as returned from a previous `HIDIOCGREPORTINFO` call.

HIDIOCGUCODE

- struct `hiddev_usage_ref` (read/write)

Returns the `usage_code` in a `hiddev_usage_ref` structure, given that given its report type, report id, field index, and index within the field have already been filled into the structure.

HIDIOCGUSAGE

- struct `hiddev_usage_ref` (read/write)

Returns the value of a usage in a `hiddev_usage_ref` structure. The usage to be retrieved can be specified as above, or the user can choose to fill in the `report_type` field and specify the `report_id` as `HID_REPORT_ID_UNKNOWN`. In this case, the `hiddev_usage_ref` will be filled in with the report and field information associated with this usage if it is found.

HIDIOCSUSAGE

- struct `hiddev_usage_ref` (write)

Sets the value of a usage in an output report. The user fills in the `hiddev_usage_ref` structure as above, but additionally fills in the `value` field.

HIDIOGCOLLECTIONINDEX

- struct hiddev_usage_ref (write)

Returns the collection index associated with this usage. This indicates where in the collection hierarchy this usage sits.

HIDIOCGFLAG

- int (read)

HIDIOCSFLAG

- int (write)

These operations respectively inspect and replace the mode flags that influence the read() call above. The flags are as follows:

HIDDEV_FLAG_UREF

- read() calls will now return struct hiddev_usage_ref instead of struct hiddev_event. This is a larger structure, but in situations where the device has more than one usage in its reports with the same usage code, this mode serves to resolve such ambiguity.

HIDDEV_FLAG_REPORT

- This flag can only be used in conjunction with HIDDEV_FLAG_UREF. With this flag set, when the device sends a report, a struct hiddev_usage_ref will be returned to read() filled in with the report_type and report_id, but with field_index set to FIELD_INDEX_NONE. This serves as additional notification when the device has sent a report.

HIDRAW - RAW ACCESS TO USB AND BLUETOOTH HUMAN INTERFACE DEVICES

The hidraw driver provides a raw interface to USB and Bluetooth Human Interface Devices (HIDs). It differs from hiddev in that reports sent and received are not parsed by the HID parser, but are sent to and received from the device unmodified.

Hidraw should be used if the userspace application knows exactly how to communicate with the hardware device, and is able to construct the HID reports manually. This is often the case when making userspace drivers for custom HID devices.

Hidraw is also useful for communicating with non-conformant HID devices which send and receive data in a way that is inconsistent with their report descriptors. Because hiddev parses reports which are sent and received through it, checking them against the device's report descriptor, such communication with these non-conformant devices is impossible using hiddev. Hidraw is the only alternative, short of writing a custom kernel driver, for these non-conformant devices.

A benefit of hidraw is that its use by userspace applications is independent of the underlying hardware type. Currently, Hidraw is implemented for USB and Bluetooth. In the future, as new hardware bus types are developed which use the HID specification, hidraw will be expanded to add support for these new bus types.

Hidraw uses a dynamic major number, meaning that udev should be relied on to create hidraw device nodes. Udev will typically create the device nodes directly under /dev (eg: /dev/hidraw0). As this location is distribution- and udev rule-dependent, applications should use libudev to locate hidraw devices attached to the system. There is a tutorial on libudev with a working example at:

<http://www.signal11.us/oss/udev/>

2.1 The HIDRAW API

2.2 read()

read() will read a queued report received from the HID device. On USB devices, the reports read using read() are the reports sent from the device on the INTERRUPT IN endpoint. By default, read() will block until there is a report available to be read. read() can be made non-blocking, by passing the O_NONBLOCK flag to open(), or by setting the O_NONBLOCK flag using fcntl().

On a device which uses numbered reports, the first byte of the returned data will be the report number; the report data follows, beginning in the second byte. For devices which do not use numbered reports, the report data will begin at the first byte.

2.3 write()

The write() function will write a report to the device. For USB devices, if the device has an INTERRUPT OUT endpoint, the report will be sent on that endpoint. If it does not, the report will be sent over the control endpoint, using a SET_REPORT transfer.

The first byte of the buffer passed to write() should be set to the report number. If the device does not use numbered reports, the first byte should be set to 0. The report data itself should begin at the second byte.

2.4 ioctl()

Hidraw supports the following ioctls:

HIDIOCGRDESCSIZE: Get Report Descriptor Size

This ioctl will get the size of the device' s report descriptor.

HIDIOCGRDESC: Get Report Descriptor

This ioctl returns the device' s report descriptor using a hidraw_report_descriptor struct. Make sure to set the size field of the hidraw_report_descriptor struct to the size returned from HIDIOCGRDESCSIZE.

HIDIOCGRAWINFO: Get Raw Info

This ioctl will return a hidraw_devinfo struct containing the bus type, the vendor ID (VID), and product ID (PID) of the device. The bus type can be one of:

- | |
|---|
| <ul style="list-style-type: none">- BUS_USB- BUS_HIL- BUS_BLUETOOTH- BUS_VIRTUAL |
|---|

which are defined in uapi/linux/input.h.

HIDIOCGRAWNAME(len): Get Raw Name

This ioctl returns a string containing the vendor and product strings of the device. The returned string is Unicode, UTF-8 encoded.

HIDIOCGRAWPHYS(len): Get Physical Address

This ioctl returns a string representing the physical address of the device. For USB devices, the string contains the physical path to the device (the USB controller, hubs, ports, etc). For Bluetooth devices, the string contains the hardware (MAC) address of the device.

HIDIOCSFEATURE(len): Send a Feature Report

This ioctl will send a feature report to the device. Per the HID specification, feature reports are always sent using the control endpoint. Set the first byte of the supplied buffer to the report number. For devices which do not use numbered reports, set the first byte to 0. The report data begins in the second byte. Make sure to set len accordingly, to one more than the length of the report (to account for the report number).

HIDIOCGFEATURE(len): Get a Feature Report

This ioctl will request a feature report from the device using the control endpoint. The first byte of the supplied buffer should be set to the report number of the requested report. For devices which do not use numbered reports, set the first byte to 0. The report will be returned starting at the first byte of the buffer (ie: the report number is not returned).

2.5 Example

In samples/, find hid-example.c, which shows examples of read(), write(), and all the ioctls for hidraw. The code may be used by anyone for any purpose, and can serve as a starting point for developing applications using hidraw.

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HID SENSORS FRAMEWORK

HID sensor framework provides necessary interfaces to implement sensor drivers, which are connected to a sensor hub. The sensor hub is a HID device and it provides a report descriptor conforming to HID 1.12 sensor usage tables.

Description from the HID 1.12 “HID Sensor Usages” specification: “Standardization of HID usages for sensors would allow (but not require) sensor hardware vendors to provide a consistent Plug And Play interface at the USB boundary, thereby enabling some operating systems to incorporate common device drivers that could be reused between vendors, alleviating any need for the vendors to provide the drivers themselves.”

This specification describes many usage IDs, which describe the type of sensor and also the individual data fields. Each sensor can have variable number of data fields. The length and order is specified in the report descriptor. For example a part of report descriptor can look like:

```
INPUT(1)[INPUT]
..
  Field(2)
    Physical(0020.0073)
    Usage(1)
      0020.045f
    Logical Minimum(-32767)
    Logical Maximum(32767)
    Report Size(8)
    Report Count(1)
    Report Offset(16)
    Flags(Variable Absolute)
..
..
```

The report is indicating “sensor page (0x20)” contains an accelerometer-3D (0x73). This accelerometer-3D has some fields. Here for example field 2 is motion intensity (0x045f) with a logical minimum value of -32767 and logical maximum of 32767. The order of fields and length of each field is important as the input event raw data will use this format.

3.1 Implementation

This specification defines many different types of sensors with different sets of data fields. It is difficult to have a common input event to user space applications, for different sensors. For example an accelerometer can send X,Y and Z data, whereas an ambient light sensor can send illumination data. So the implementation has two parts:

- Core hid driver
- Individual sensor processing part (sensor drivers)

3.1.1 Core driver

The core driver registers (hid-sensor-hub) registers as a HID driver. It parses report descriptors and identifies all the sensors present. It adds an MFD device with name HID-SENSOR-xxxx (where xxxx is usage id from the specification).

For example:

HID-SENSOR-200073 is registered for an Accelerometer 3D driver.

So if any driver with this name is inserted, then the probe routine for that function will be called. So an accelerometer processing driver can register with this name and will be probed if there is an accelerometer-3D detected.

The core driver provides a set of APIs which can be used by the processing drivers to register and get events for that usage id. Also it provides parsing functions, which get and set each input/feature/output report.

3.1.2 Individual sensor processing part (sensor drivers)

The processing driver will use an interface provided by the core driver to parse the report and get the indexes of the fields and also can get events. This driver can use IIO interface to use the standard ABI defined for a type of sensor.

3.2 Core driver Interface

Callback structure:

```
Each processing driver can use this structure to set some callbacks.
    int (*suspend)(..): Callback when HID suspend is received
    int (*resume)(..): Callback when HID resume is received
    int (*capture_sample)(..): Capture a sample for one of its data_
↳fields
    int (*send_event)(..): One complete event is received which can have
                           multiple data fields.
```

Registration functions:


```
int sensor_hub_register_callback(struct hid_sensor_hub_device *hsdev,
                               u32 usage_id,
                               struct hid_sensor_hub_callbacks *usage_callback):
```

Registers callbacks for an usage id. The callback functions are not allowed to sleep:

```
int sensor_hub_remove_callback(struct hid_sensor_hub_device *hsdev,
                              u32 usage_id):
```

Removes callbacks for an usage id.

Parsing function:

```
int sensor_hub_input_get_attribute_info(struct hid_sensor_hub_device *
↳hsdev,
                                     u8 type,
                                     u32 usage_id, u32 attr_usage_id,
                                     struct hid_sensor_hub_attribute_info *info);
```

A processing driver can look for some field of interest and check if it exists in a report descriptor. If it exists it will store necessary information so that fields can be set or get individually. These indexes avoid searching every time and getting field index to get or set.

Set Feature report:

```
int sensor_hub_set_feature(struct hid_sensor_hub_device *hsdev, u32 report_
↳id,
                          u32 field_index, s32 value);
```

This interface is used to set a value for a field in feature report. For example if there is a field `report_interval`, which is parsed by a call to `sensor_hub_input_get_attribute_info` before, then it can directly set that individual field:

```
int sensor_hub_get_feature(struct hid_sensor_hub_device *hsdev, u32 report_
↳id,
                          u32 field_index, s32 *value);
```

This interface is used to get a value for a field in input report. For example if there is a field `report_interval`, which is parsed by a call to `sensor_hub_input_get_attribute_info` before, then it can directly get that individual field value:

```
int sensor_hub_input_attr_get_raw_value(struct hid_sensor_hub_device *
↳hsdev,
                                       u32 usage_id,
                                       u32 attr_usage_id, u32 report_id);
```

This is used to get a particular field value through input reports. For example accelerometer wants to poll X axis value, then it can call this function with the usage id of X axis. HID sensors can provide events, so this is not necessary to poll for any field. If there is some new sample, the core driver will call registered callback function to process the sample.

3.2.1 HID Custom and generic Sensors

HID Sensor specification defines two special sensor usage types. Since they don't represent a standard sensor, it is not possible to define using Linux IIO type interfaces. The purpose of these sensors is to extend the functionality or provide a way to obfuscate the data being communicated by a sensor. Without knowing the mapping between the data and its encapsulated form, it is difficult for an application/driver to determine what data is being communicated by the sensor. This allows some differentiating use cases, where vendor can provide applications. Some common use cases are debug other sensors or to provide some events like keyboard attached/detached or lid open/close.

To allow application to utilize these sensors, here they are exported uses sysfs attribute groups, attributes and misc device interface.

An example of this representation on sysfs:

```
/sys/devices/pci0000:00/INT33C2:00/i2c-0/i2c-INT33D1:00/0018:8086:09FA.  
↳0001/HID-SENSOR-2000e1.6.auto$ tree -R
```

```
.  
├── enable_sensor  
│   ├── feature-0-200316  
│   │   ├── feature-0-200316-maximum  
│   │   ├── feature-0-200316-minimum  
│   │   ├── feature-0-200316-name  
│   │   ├── feature-0-200316-size  
│   │   ├── feature-0-200316-unit-expo  
│   │   ├── feature-0-200316-units  
│   │   └── feature-0-200316-value  
│   ├── feature-1-200201  
│   │   ├── feature-1-200201-maximum  
│   │   ├── feature-1-200201-minimum  
│   │   ├── feature-1-200201-name  
│   │   ├── feature-1-200201-size  
│   │   ├── feature-1-200201-unit-expo  
│   │   ├── feature-1-200201-units  
│   │   └── feature-1-200201-value  
│   ├── input-0-200201  
│   │   ├── input-0-200201-maximum  
│   │   ├── input-0-200201-minimum  
│   │   ├── input-0-200201-name  
│   │   ├── input-0-200201-size  
│   │   ├── input-0-200201-unit-expo  
│   │   ├── input-0-200201-units  
│   │   └── input-0-200201-value  
│   └── input-1-200202  
│       ├── input-1-200202-maximum  
│       ├── input-1-200202-minimum  
│       ├── input-1-200202-name  
│       └── input-1-200202-size
```

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```
| | | — input-1-200202-unit-expo  
| | | — input-1-200202-units  
| | | — input-1-200202-value
```

Here there is a custom sensors with four fields, two feature and two inputs. Each field is represented by a set of attributes. All fields except the “value” are read only. The value field is a RW field.

Example:

```
/sys/bus/platform/devices/HID-SENSOR-2000e1.6.auto/feature-0-200316$ grep -  
↪ r . *  
feature-0-200316-maximum:6  
feature-0-200316-minimum:0  
feature-0-200316-name:property-reporting-state  
feature-0-200316-size:1  
feature-0-200316-unit-expo:0  
feature-0-200316-units:25  
feature-0-200316-value:1
```

How to enable such sensor?

By default sensor can be power gated. To enable sysfs attribute “enable” can be used:

```
$ echo 1 > enable_sensor
```

Once enabled and powered on, sensor can report value using HID reports. These reports are pushed using misc device interface in a FIFO order:

```
/dev$ tree | grep HID-SENSOR-2000e1.6.auto  
|   |   | — 10:53 -> ../HID-SENSOR-2000e1.6.auto  
|   | — HID-SENSOR-2000e1.6.auto
```

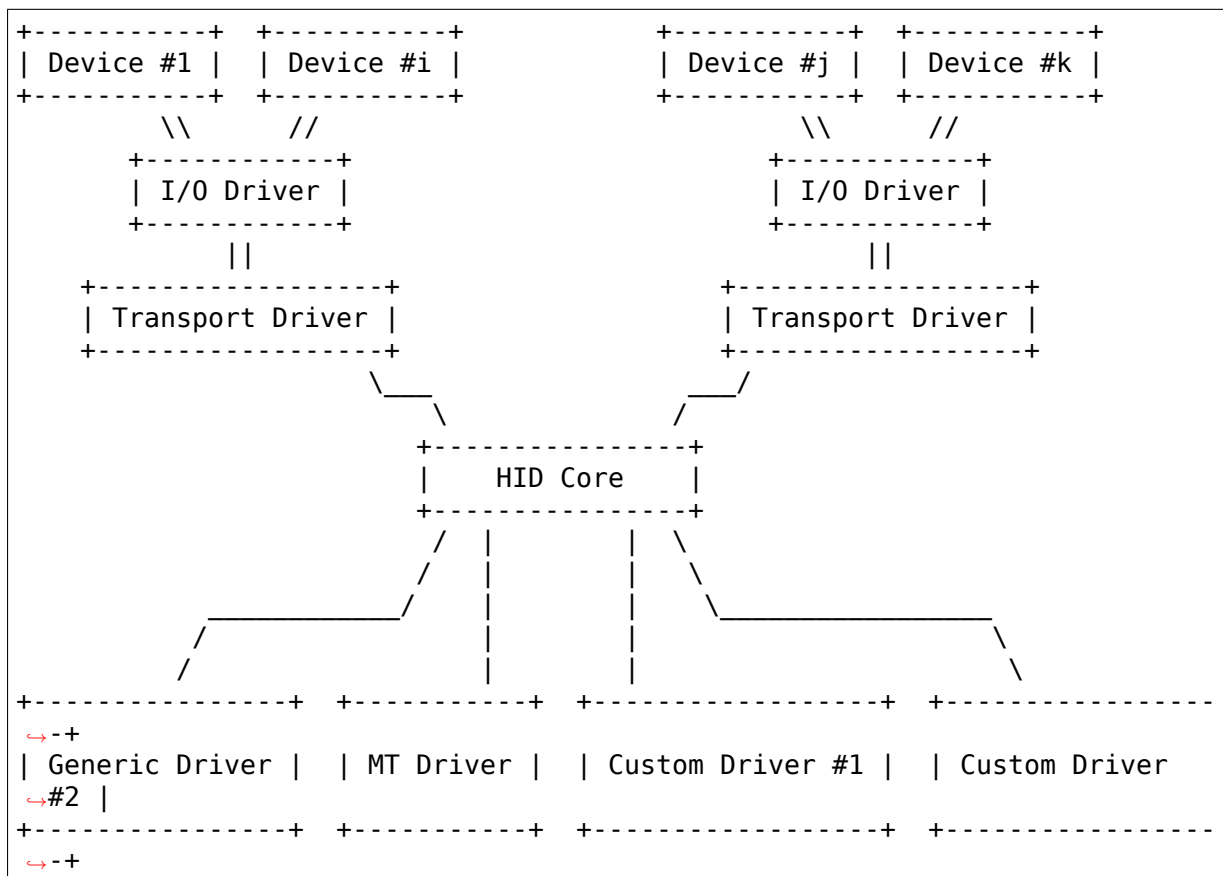
Each reports can be of variable length preceded by a header. This header consist of a 32 bit usage id, 64 bit time stamp and 32 bit length field of raw data.

HID I/O TRANSPORT DRIVERS

The HID subsystem is independent of the underlying transport driver. Initially, only USB was supported, but other specifications adopted the HID design and provided new transport drivers. The kernel includes at least support for USB, Bluetooth, I2C and user-space I/O drivers.

4.1 1) HID Bus

The HID subsystem is designed as a bus. Any I/O subsystem may provide HID devices and register them with the HID bus. HID core then loads generic device drivers on top of it. The transport drivers are responsible of raw data transport and device setup/management. HID core is responsible of report-parsing, report interpretation and the user-space API. Device specifics and quirks are handled by all layers depending on the quirk.



Example Drivers:

- I/O: USB, I2C, Bluetooth-l2cap
- Transport: USB-HID, I2C-HID, BT-HIDP

Everything below “HID Core” is simplified in this graph as it is only of interest to HID device drivers. Transport drivers do not need to know the specifics.

4.1.1 1.1) Device Setup

I/O drivers normally provide hotplug detection or device enumeration APIs to the transport drivers. Transport drivers use this to find any suitable HID device. They allocate HID device objects and register them with HID core. Transport drivers are not required to register themselves with HID core. HID core is never aware of which transport drivers are available and is not interested in it. It is only interested in devices.

Transport drivers attach a constant “struct hid_ll_driver” object with each device. Once a device is registered with HID core, the callbacks provided via this struct are used by HID core to communicate with the device.

Transport drivers are responsible of detecting device failures and unplugging. HID core will operate a device as long as it is registered regardless of any device failures. Once transport drivers detect unplug or failure events, they must unregister the device from HID core and HID core will stop using the provided callbacks.

4.1.2 1.2) Transport Driver Requirements

The terms “asynchronous” and “synchronous” in this document describe the transmission behavior regarding acknowledgements. An asynchronous channel must not perform any synchronous operations like waiting for acknowledgements or verifications. Generally, HID calls operating on asynchronous channels must be running in atomic-context just fine. On the other hand, synchronous channels can be implemented by the transport driver in whatever way they like. They might just be the same as asynchronous channels, but they can also provide acknowledgement reports, automatic retransmission on failure, etc. in a blocking manner. If such functionality is required on asynchronous channels, a transport-driver must implement that via its own worker threads.

HID core requires transport drivers to follow a given design. A Transport driver must provide two bi-directional I/O channels to each HID device. These channels must not necessarily be bi-directional in the hardware itself. A transport driver might just provide 4 uni-directional channels. Or it might multiplex all four on a single physical channel. However, in this document we will describe them as two bi-directional channels as they have several properties in common.

- Interrupt Channel (intr): The intr channel is used for asynchronous data reports. No management commands or data acknowledgements are sent on this channel. Any unrequested incoming or outgoing data report must be sent on this channel and is never acknowledged by the remote side. Devices usually send their input events on this channel. Outgoing events are normally not send via intr, except if high throughput is required.

- **Control Channel (ctrl):** The ctrl channel is used for synchronous requests and device management. Unrequested data input events must not be sent on this channel and are normally ignored. Instead, devices only send management events or answers to host requests on this channel. The control-channel is used for direct blocking queries to the device independent of any events on the intr-channel. Outgoing reports are usually sent on the ctrl channel via synchronous SET_REPORT requests.

Communication between devices and HID core is mostly done via HID reports. A report can be of one of three types:

- **INPUT Report:** Input reports provide data from device to host. This data may include button events, axis events, battery status or more. This data is generated by the device and sent to the host with or without requiring explicit requests. Devices can choose to send data continuously or only on change.
- **OUTPUT Report:** Output reports change device states. They are sent from host to device and may include LED requests, rumble requests or more. Output reports are never sent from device to host, but a host can retrieve their current state. Hosts may choose to send output reports either continuously or only on change.
- **FEATURE Report:** Feature reports are used for specific static device features and never reported spontaneously. A host can read and/or write them to access data like battery-state or device-settings. Feature reports are never sent without requests. A host must explicitly set or retrieve a feature report. This also means, feature reports are never sent on the intr channel as this channel is asynchronous.

INPUT and OUTPUT reports can be sent as pure data reports on the intr channel. For INPUT reports this is the usual operational mode. But for OUTPUT reports, this is rarely done as OUTPUT reports are normally quite scarce. But devices are free to make excessive use of asynchronous OUTPUT reports (for instance, custom HID audio speakers make great use of it).

Plain reports must not be sent on the ctrl channel, though. Instead, the ctrl channel provides synchronous GET/SET_REPORT requests. Plain reports are only allowed on the intr channel and are the only means of data there.

- **GET_REPORT:** A GET_REPORT request has a report ID as payload and is sent from host to device. The device must answer with a data report for the requested report ID on the ctrl channel as a synchronous acknowledgement. Only one GET_REPORT request can be pending for each device. This restriction is enforced by HID core as several transport drivers don't allow multiple simultaneous GET_REPORT requests. Note that data reports which are sent as answer to a GET_REPORT request are not handled as generic device events. That is, if a device does not operate in continuous data reporting mode, an answer to GET_REPORT does not replace the raw data report on the intr channel on state change. GET_REPORT is only used by custom HID device drivers to query device state. Normally, HID core caches any device state so this request is not necessary on devices that follow the HID specs except during device initialization to retrieve the current state. GET_REPORT requests can be sent for any of the 3 report types and shall return the current report state of the device. However, OUTPUT reports as payload may be blocked by the underlying transport driver if the specification does not allow

them.

- **SET_REPORT**: A SET_REPORT request has a report ID plus data as payload. It is sent from host to device and a device must update its current report state according to the given data. Any of the 3 report types can be used. However, INPUT reports as payload might be blocked by the underlying transport driver if the specification does not allow them. A device must answer with a synchronous acknowledgement. However, HID core does not require transport drivers to forward this acknowledgement to HID core. Same as for GET_REPORT, only one SET_REPORT can be pending at a time. This restriction is enforced by HID core as some transport drivers do not support multiple synchronous SET_REPORT requests.

Other ctrl-channel requests are supported by USB-HID but are not available (or deprecated) in most other transport level specifications:

- **GET/SET_IDLE**: Only used by USB-HID and I2C-HID.
- **GET/SET_PROTOCOL**: Not used by HID core.
- **RESET**: Used by I2C-HID, not hooked up in HID core.
- **SET_POWER**: Used by I2C-HID, not hooked up in HID core.

4.2 2) HID API

4.2.1 2.1) Initialization

Transport drivers normally use the following procedure to register a new device with HID core:

```
struct hid_device *hid;
int ret;

hid = hid_allocate_device();
if (IS_ERR(hid)) {
    ret = PTR_ERR(hid);
    goto err_<...>;
}

strncpy(hid->name, <device-name-src>, sizeof(hid->name));
strncpy(hid->phys, <device-phys-src>, sizeof(hid->phys));
strncpy(hid->uniq, <device-uniq-src>, sizeof(hid->uniq));

hid->ll_driver = &custom_ll_driver;
hid->bus = <device-bus>;
hid->vendor = <device-vendor>;
hid->product = <device-product>;
hid->version = <device-version>;
hid->country = <device-country>;
hid->dev.parent = <pointer-to-parent-device>;
hid->driver_data = <transport-driver-data-field>;

ret = hid_add_device(hid);
```

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```
if (ret)
    goto err_<...>;
```

Once `hid_add_device()` is entered, HID core might use the callbacks provided in “`custom_ll_driver`”. Note that fields like “`country`” can be ignored by underlying transport-drivers if not supported.

To unregister a device, use:

```
hid_destroy_device(hid);
```

Once `hid_destroy_device()` returns, HID core will no longer make use of any driver callbacks.

4.2.2 2.2) hid_ll_driver operations

The available HID callbacks are:

```
int (*start) (struct hid_device *hdev)
```

Called from HID device drivers once they want to use the device. Transport drivers can choose to setup their device in this callback. However, normally devices are already set up before transport drivers register them to HID core so this is mostly only used by USB-HID.

```
void (*stop) (struct hid_device *hdev)
```

Called from HID device drivers once they are done with a device. Transport drivers can free any buffers and deinitialize the device. But note that `->start()` might be called again if another HID device driver is loaded on the device.

Transport drivers are free to ignore it and deinitialize devices after they destroyed them via `hid_destroy_device()`.

```
int (*open) (struct hid_device *hdev)
```

Called from HID device drivers once they are interested in data reports. Usually, while user-space didn't open any input API/etc., device drivers are not interested in device data and transport drivers can put devices asleep. However, once `->open()` is called, transport drivers must be ready for I/O. `->open()` calls are nested for each client that opens the HID device.

```
void (*close) (struct hid_device *hdev)
```

Called from HID device drivers after `->open()` was called but they are no longer interested in device reports. (Usually if user-space closed any input devices of the driver).

Transport drivers can put devices asleep and terminate any I/O of all `->open()` calls have been followed by a `->close()` call. However, `->start()`

may be called again if the device driver is interested in input reports again.

```
int (*parse) (struct hid_device *hdev)
```

Called once during device setup after `->start()` has been called. Transport drivers must read the HID report-descriptor from the device and tell HID core about it via `hid_parse_report()`.

```
int (*power) (struct hid_device *hdev, int level)
```

Called by HID core to give PM hints to transport drivers. Usually this is analogical to the `->open()` and `->close()` hints and redundant.

```
void (*request) (struct hid_device *hdev, struct hid_report_  
↪*report,  
                int reqtype)
```

Send an HID request on the ctrl channel. “report” contains the report that should be sent and “reqtype” the request type. Request-type can be `HID_REQ_SET_REPORT` or `HID_REQ_GET_REPORT`.

This callback is optional. If not provided, HID core will assemble a raw report following the HID specs and send it via the `->raw_request()` callback. The transport driver is free to implement this asynchronously.

```
int (*wait) (struct hid_device *hdev)
```

Used by HID core before calling `->request()` again. A transport driver can use it to wait for any pending requests to complete if only one request is allowed at a time.

```
int (*raw_request) (struct hid_device *hdev, unsigned char_  
↪reportnum,  
                  __u8 *buf, size_t count, unsigned char rtype,  
                  int reqtype)
```

Same as `->request()` but provides the report as raw buffer. This request shall be synchronous. A transport driver must not use `->wait()` to complete such requests. This request is mandatory and hid core will reject the device if it is missing.

```
int (*output_report) (struct hid_device *hdev, __u8 *buf, size_t_  
↪len)
```

Send raw output report via intr channel. Used by some HID device drivers which require high throughput for outgoing requests on the intr channel. This must not cause `SET_REPORT` calls! This must be implemented as asynchronous output report on the intr channel!

```
int (*idle) (struct hid_device *hdev, int report, int idle, int_  
↪reqtype)
```

Perform `SET/GET_IDLE` request. Only used by USB-HID, do not implement!

4.2.3 2.3) Data Path

Transport drivers are responsible of reading data from I/O devices. They must handle any I/O-related state-tracking themselves. HID core does not implement protocol handshakes or other management commands which can be required by the given HID transport specification.

Every raw data packet read from a device must be fed into HID core via `hid_input_report()`. You must specify the channel-type (`intr` or `ctrl`) and report type (`input/output/feature`). Under normal conditions, only input reports are provided via this API.

Responses to `GET_REPORT` requests via `->request()` must also be provided via this API. Responses to `->raw_request()` are synchronous and must be intercepted by the transport driver and not passed to `hid_input_report()`. Acknowledgements to `SET_REPORT` requests are not of interest to HID core.

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UHID - USER-SPACE I/O DRIVER SUPPORT FOR HID SUBSYSTEM

UHID allows user-space to implement HID transport drivers. Please see `hid-transport.txt` for an introduction into HID transport drivers. This document relies heavily on the definitions declared there.

With UHID, a user-space transport driver can create kernel hid-devices for each device connected to the user-space controlled bus. The UHID API defines the I/O events provided from the kernel to user-space and vice versa.

There is an example user-space application in `./samples/uhid/uhid-example.c`

5.1 The UHID API

UHID is accessed through a character misc-device. The minor-number is allocated dynamically so you need to rely on `udev` (or similar) to create the device node. This is `/dev/uhid` by default.

If a new device is detected by your HID I/O Driver and you want to register this device with the HID subsystem, then you need to open `/dev/uhid` once for each device you want to register. All further communication is done by `read()`'ing or `write()`'ing "struct `uhid_event`" objects. Non-blocking operations are supported by setting `O_NONBLOCK`:

```
struct uhid_event {
    __u32 type;
    union {
        struct uhid_create2_req create2;
        struct uhid_output_req output;
        struct uhid_input2_req input2;
        ...
    } u;
};
```

The "type" field contains the ID of the event. Depending on the ID different payloads are sent. You must not split a single event across multiple `read()`'s or multiple `write()`'s. A single event must always be sent as a whole. Furthermore, only a single event can be sent per `read()` or `write()`. Pending data is ignored. If you want to handle multiple events in a single syscall, then use vectored I/O with `readv()/writev()`. The "type" field defines the payload. For each type, there is a

payload-structure available in the union “u” (except for empty payloads). This payload contains management and/or device data.

The first thing you should do is sending an `UHID_CREATE2` event. This will register the device. UHID will respond with an `UHID_START` event. You can now start sending data to and reading data from UHID. However, unless UHID sends the `UHID_OPEN` event, the internally attached HID Device Driver has no user attached. That is, you might put your device asleep unless you receive the `UHID_OPEN` event. If you receive the `UHID_OPEN` event, you should start I/O. If the last user closes the HID device, you will receive an `UHID_CLOSE` event. This may be followed by an `UHID_OPEN` event again and so on. There is no need to perform reference-counting in user-space. That is, you will never receive multiple `UHID_OPEN` events without an `UHID_CLOSE` event. The HID subsystem performs ref-counting for you. You may decide to ignore `UHID_OPEN`/`UHID_CLOSE`, though. I/O is allowed even though the device may have no users.

If you want to send data on the interrupt channel to the HID subsystem, you send an `HID_INPUT2` event with your raw data payload. If the kernel wants to send data on the interrupt channel to the device, you will read an `UHID_OUTPUT` event. Data requests on the control channel are currently limited to `GET_REPORT` and `SET_REPORT` (no other data reports on the control channel are defined so far). Those requests are always synchronous. That means, the kernel sends `UHID_GET_REPORT` and `UHID_SET_REPORT` events and requires you to forward them to the device on the control channel. Once the device responds, you must forward the response via `UHID_GET_REPORT_REPLY` and `UHID_SET_REPORT_REPLY` to the kernel. The kernel blocks internal driver-execution during such round-trips (times out after a hard-coded period).

If your device disconnects, you should send an `UHID_DESTROY` event. This will unregister the device. You can now send `UHID_CREATE2` again to register a new device. If you `close()` the `fd`, the device is automatically unregistered and destroyed internally.

5.2 write()

`write()` allows you to modify the state of the device and feed input data into the kernel. The kernel will parse the event immediately and if the event ID is not supported, it will return `-EOPNOTSUPP`. If the payload is invalid, then `-EINVAL` is returned, otherwise, the amount of data that was read is returned and the request was handled successfully. `O_NONBLOCK` does not affect `write()` as writes are always handled immediately in a non-blocking fashion. Future requests might make use of `O_NONBLOCK`, though.

UHID_CREATE2: This creates the internal HID device. No I/O is possible until you send this event to the kernel. The payload is of type `struct uhid_create2_req` and contains information about your device. You can start I/O now.

UHID_DESTROY: This destroys the internal HID device. No further I/O will be accepted. There may still be pending messages that you can receive with `read()` but no further `UHID_INPUT` events can be sent to the kernel. You can create a new device by sending `UHID_CREATE2` again. There is no need to

reopen the character device.

UHID_INPUT2: You must send UHID_CREATE2 before sending input to the kernel! This event contains a data-payload. This is the raw data that you read from your device on the interrupt channel. The kernel will parse the HID reports.

UHID_GET_REPORT_REPLY: If you receive a UHID_GET_REPORT request you must answer with this request. You must copy the “id” field from the request into the answer. Set the “err” field to 0 if no error occurred or to EIO if an I/O error occurred. If “err” is 0 then you should fill the buffer of the answer with the results of the GET_REPORT request and set “size” correspondingly.

UHID_SET_REPORT_REPLY: This is the SET_REPORT equivalent of UHID_GET_REPORT_REPLY. Unlike GET_REPORT, SET_REPORT never returns a data buffer, therefore, it’s sufficient to set the “id” and “err” fields correctly.

5.3 read()

read() will return a queued output report. No reaction is required to any of them but you should handle them according to your needs.

UHID_START: This is sent when the HID device is started. Consider this as an answer to UHID_CREATE2. This is always the first event that is sent. Note that this event might not be available immediately after write(UHID_CREATE2) returns. Device drivers might required delayed setups. This event contains a payload of type uhid_start_req. The “dev_flags” field describes special behaviors of a device. The following flags are defined:

- UHID_DEV_NUMBERED_FEATURE_REPORTS
- UHID_DEV_NUMBERED_OUTPUT_REPORTS
- UHID_DEV_NUMBERED_INPUT_REPORTS

Each of these flags defines whether a given report-type uses numbered reports. If numbered reports are used for a type, all messages from the kernel already have the report-number as prefix. Otherwise, no prefix is added by the kernel. For messages sent by user-space to the kernel, you must adjust the prefixes according to these flags.

UHID_STOP: This is sent when the HID device is stopped. Consider this as an answer to UHID_DESTROY.

If you didn’t destroy your device via UHID_DESTROY, but the kernel sends an UHID_STOP event, this should usually be ignored. It means that the kernel reloaded/changed the device driver loaded on your HID device (or some other maintenance actions happened).

You can usually ignored any UHID_STOP events safely.

UHID_OPEN: This is sent when the HID device is opened. That is, the data that the HID device provides is read by some other process. You may ignore this event but it is useful for power-management. As long as you haven’t received

this event there is actually no other process that reads your data so there is no need to send UHID_INPUT2 events to the kernel.

UHID_CLOSE: This is sent when there are no more processes which read the HID data. It is the counterpart of UHID_OPEN and you may as well ignore this event.

UHID_OUTPUT: This is sent if the HID device driver wants to send raw data to the I/O device on the interrupt channel. You should read the payload and forward it to the device. The payload is of type “struct uhid_output_req”. This may be received even though you haven’t received UHID_OPEN, yet.

UHID_GET_REPORT: This event is sent if the kernel driver wants to perform a GET_REPORT request on the control channel as described in the HID specs. The report-type and report-number are available in the payload. The kernel serializes GET_REPORT requests so there will never be two in parallel. However, if you fail to respond with a UHID_GET_REPORT_REPLY, the request might silently time out. Once you read a GET_REPORT request, you shall forward it to the hid device and remember the “id” field in the payload. Once your hid device responds to the GET_REPORT (or if it fails), you must send a UHID_GET_REPORT_REPLY to the kernel with the exact same “id” as in the request. If the request already timed out, the kernel will ignore the response silently. The “id” field is never re-used, so conflicts cannot happen.

UHID_SET_REPORT: This is the SET_REPORT equivalent of UHID_GET_REPORT. On receipt, you shall send a SET_REPORT request to your hid device. Once it replies, you must tell the kernel about it via UHID_SET_REPORT_REPLY. The same restrictions as for UHID_GET_REPORT apply.

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ALPS HID TOUCHPAD PROTOCOL

6.1 Introduction

Currently ALPS HID driver supports U1 Touchpad device.

U1 device basic information.

Vender ID	0x044E
Product ID	0x120B
Version ID	0x0121

6.2 HID Descriptor

Byte	Field	Value	Notes
0	wHIDDescLength	001E	Length of HID Descriptor : 30 bytes
2	bcdVersion	0100	Compliant with Version 1.00
4	wReportDescLength	00B2	Report Descriptor is 178 Bytes (0x00B2)
6	wReportDescRegister	0002	Identifier to read Report Descriptor
8	wInputRegister	0003	Identifier to read Input Report
10	wMaxInputLength	0053	Input Report is 80 Bytes + 2
12	wOutputRegister	0000	Identifier to read Output Report
14	wMaxOutputLength	0000	No Output Reports
16	wCommandRegister	0005	Identifier for Command Register
18	wDataRegister	0006	Identifier for Data Register
20	wVendorID	044E	Vendor ID 0x044E
22	wProductID	120B	Product ID 0x120B
24	wVersionID	0121	Version 01.21
26	RESERVED	0000	RESERVED

6.3 Report ID

ReportID-1	(Input Reports)	(HIDUsage-Mouse) for TP&SP
ReportID-2	(Input Reports)	(HIDUsage-keyboard) for TP
ReportID-3	(Input Reports)	(Vendor Usage: Max 10 finger data) for TP
ReportID-4	(Input Reports)	(Vendor Usage: ON bit data) for GP
ReportID-5	(Feature Reports)	Feature Reports
ReportID-6	(Input Reports)	(Vendor Usage: StickPointer data) for SP
ReportID-7	(Feature Reports)	Flash update (Bootloader)

6.4 Data pattern

Case1	ReportID_1	TP/SP	Relative/Relative
Case2	ReportID_3 ReportID_6	TP SP	Absolute Absolute

6.5 Command Read/Write

To read/write to RAM, need to send a commands to the device.

The command format is as below.

DataByte(SET_REPORT)

Byte1	Command Byte
Byte2	Address - Byte 0 (LSB)
Byte3	Address - Byte 1
Byte4	Address - Byte 2
Byte5	Address - Byte 3 (MSB)
Byte6	Value Byte
Byte7	Checksum

Command Byte is read=0xD1/write=0xD2 .

Address is read/write RAM address.

Value Byte is writing data when you send the write commands.

When you read RAM, there is no meaning.

DataByte(GET_REPORT)

Byte1	Response Byte
Byte2	Address - Byte 0 (LSB)
Byte3	Address - Byte 1
Byte4	Address - Byte 2
Byte5	Address - Byte 3 (MSB)
Byte6	Value Byte
Byte7	Checksum

Read value is stored in Value Byte.

Packet Format Touchpad data byte _____

.	b7	b6	b5	b4	b3	b2	b1	b0
1	0	0	SW6	SW5	SW4	SW3	SW2	SW1
2	0	0	0	Fcv	Fn3	Fn2	Fn1	Fn0
3	Xa0_7	Xa0_6	Xa0_5	Xa0_4	Xa0_3	Xa0_2	Xa0_1	Xa0_0
4	Xa0_15	Xa0_14	Xa0_13	Xa0_12	Xa0_11	Xa0_10	Xa0_9	Xa0_8
5	Ya0_7	Ya0_6	Ya0_5	Ya0_4	Ya0_3	Ya0_2	Ya0_1	Ya0_0
6	Ya0_15	Ya0_14	Ya0_13	Ya0_12	Ya0_11	Ya0_10	Ya0_9	Ya0_8
7	LFB0	Zs0_6	Zs0_5	Zs0_4	Zs0_3	Zs0_2	Zs0_1	Zs0_0
8	Xa1_7	Xa1_6	Xa1_5	Xa1_4	Xa1_3	Xa1_2	Xa1_1	Xa1_0
9	Xa1_15	Xa1_14	Xa1_13	Xa1_12	Xa1_11	Xa1_10	Xa1_9	Xa1_8
10	Ya1_7	Ya1_6	Ya1_5	Ya1_4	Ya1_3	Ya1_2	Ya1_1	Ya1_0
11	Ya1_15	Ya1_14	Ya1_13	Ya1_12	Ya1_11	Ya1_10	Ya1_9	Ya1_8
12	LFB1	Zs1_6	Zs1_5	Zs1_4	Zs1_3	Zs1_2	Zs1_1	Zs1_0
13	Xa2_7	Xa2_6	Xa2_5	Xa2_4	Xa2_3	Xa2_2	Xa2_1	Xa2_0
14	Xa2_15	Xa2_14	Xa2_13	Xa2_12	Xa2_11	Xa2_10	Xa2_9	Xa2_8
15	Ya2_7	Ya2_6	Ya2_5	Ya2_4	Ya2_3	Ya2_2	Ya2_1	Ya2_0
16	Ya2_15	Ya2_14	Ya2_13	Ya2_12	Ya2_11	Ya2_10	Ya2_9	Ya2_8
17	LFB2	Zs2_6	Zs2_5	Zs2_4	Zs2_3	Zs2_2	Zs2_1	Zs2_0
18	Xa3_7	Xa3_6	Xa3_5	Xa3_4	Xa3_3	Xa3_2	Xa3_1	Xa3_0
19	Xa3_15	Xa3_14	Xa3_13	Xa3_12	Xa3_11	Xa3_10	Xa3_9	Xa3_8
20	Ya3_7	Ya3_6	Ya3_5	Ya3_4	Ya3_3	Ya3_2	Ya3_1	Ya3_0
21	Ya3_15	Ya3_14	Ya3_13	Ya3_12	Ya3_11	Ya3_10	Ya3_9	Ya3_8
22	LFB3	Zs3_6	Zs3_5	Zs3_4	Zs3_3	Zs3_2	Zs3_1	Zs3_0
23	Xa4_7	Xa4_6	Xa4_5	Xa4_4	Xa4_3	Xa4_2	Xa4_1	Xa4_0
24	Xa4_15	Xa4_14	Xa4_13	Xa4_12	Xa4_11	Xa4_10	Xa4_9	Xa4_8
25	Ya4_7	Ya4_6	Ya4_5	Ya4_4	Ya4_3	Ya4_2	Ya4_1	Ya4_0
26	Ya4_15	Ya4_14	Ya4_13	Ya4_12	Ya4_11	Ya4_10	Ya4_9	Ya4_8
27	LFB4	Zs4_6	Zs4_5	Zs4_4	Zs4_3	Zs4_2	Zs4_1	Zs4_0

SW1-SW6: SW ON/OFF status

Xan_15-0(16bit): X Absolute data of the “n” th finger

Yan_15-0(16bit): Y Absolute data of the “n” th finger

Zsn_6-0(7bit): Operation area of the “n” th finger

6.6 StickPointer data byte

.	b7	b6	b5	b4	b3	b2	b1	b0
Byte1	1	1	1	0	1	SW3	SW2	SW1
Byte2	X7	X6	X5	X4	X3	X2	X1	X0
Byte3	X15	X14	X13	X12	X11	X10	X9	X8
Byte4	Y7	Y6	Y5	Y4	Y3	Y2	Y1	Y0
Byte5	Y15	Y14	Y13	Y12	Y11	Y10	Y9	Y8
Byte6	Z7	Z6	Z5	Z4	Z3	Z2	Z1	Z0
Byte7	T&P	Z14	Z13	Z12	Z11	Z10	Z9	Z8

SW1-SW3: SW ON/OFF status

Xn_15-0(16bit): X Absolute data

Yn_15-0(16bit): Y Absolute data

Zn_14-0(15bit): Z

INTEL INTEGRATED SENSOR HUB (ISH)

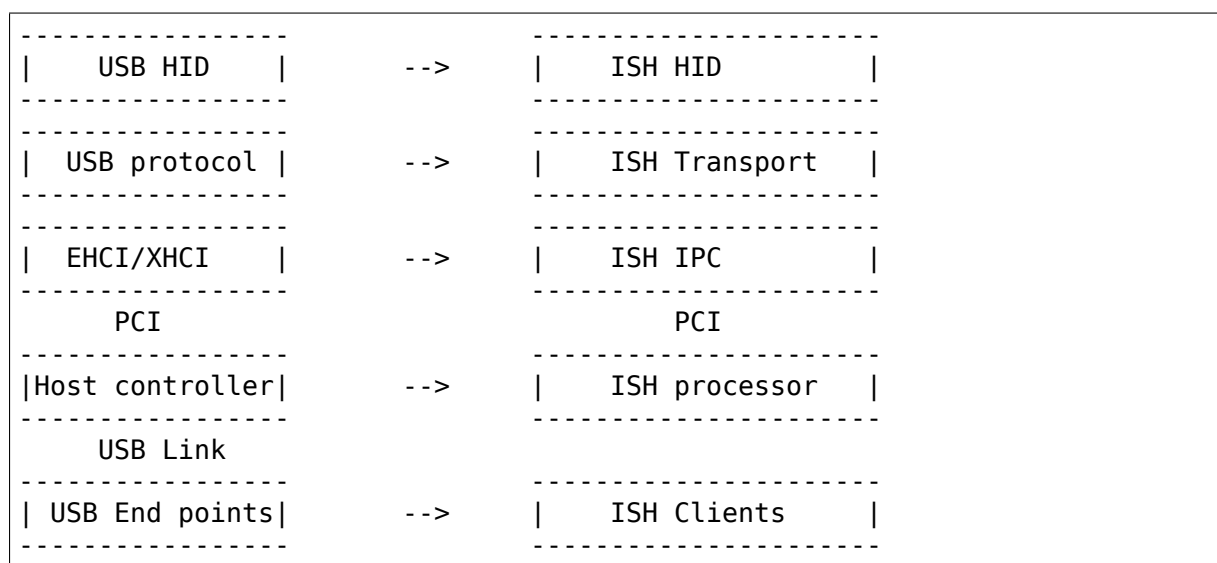
A sensor hub enables the ability to offload sensor polling and algorithm processing to a dedicated low power co-processor. This allows the core processor to go into low power modes more often, resulting in the increased battery life.

There are many vendors providing external sensor hubs conforming to HID Sensor usage tables, and used in several tablets, 2 in 1 convertible laptops and embedded products. Linux had this support since Linux 3.9.

Intel® introduced integrated sensor hubs as a part of the SoC starting from Cherry Trail and now supported on multiple generations of CPU packages. There are many commercial devices already shipped with Integrated Sensor Hubs (ISH). These ISH also comply to HID sensor specification, but the difference is the transport protocol used for communication. The current external sensor hubs mainly use HID over i2C or USB. But ISH doesn't use either i2c or USB.

7.1 1. Overview

Using a analogy with a usbhid implementation, the ISH follows a similar model for a very high speed communication:



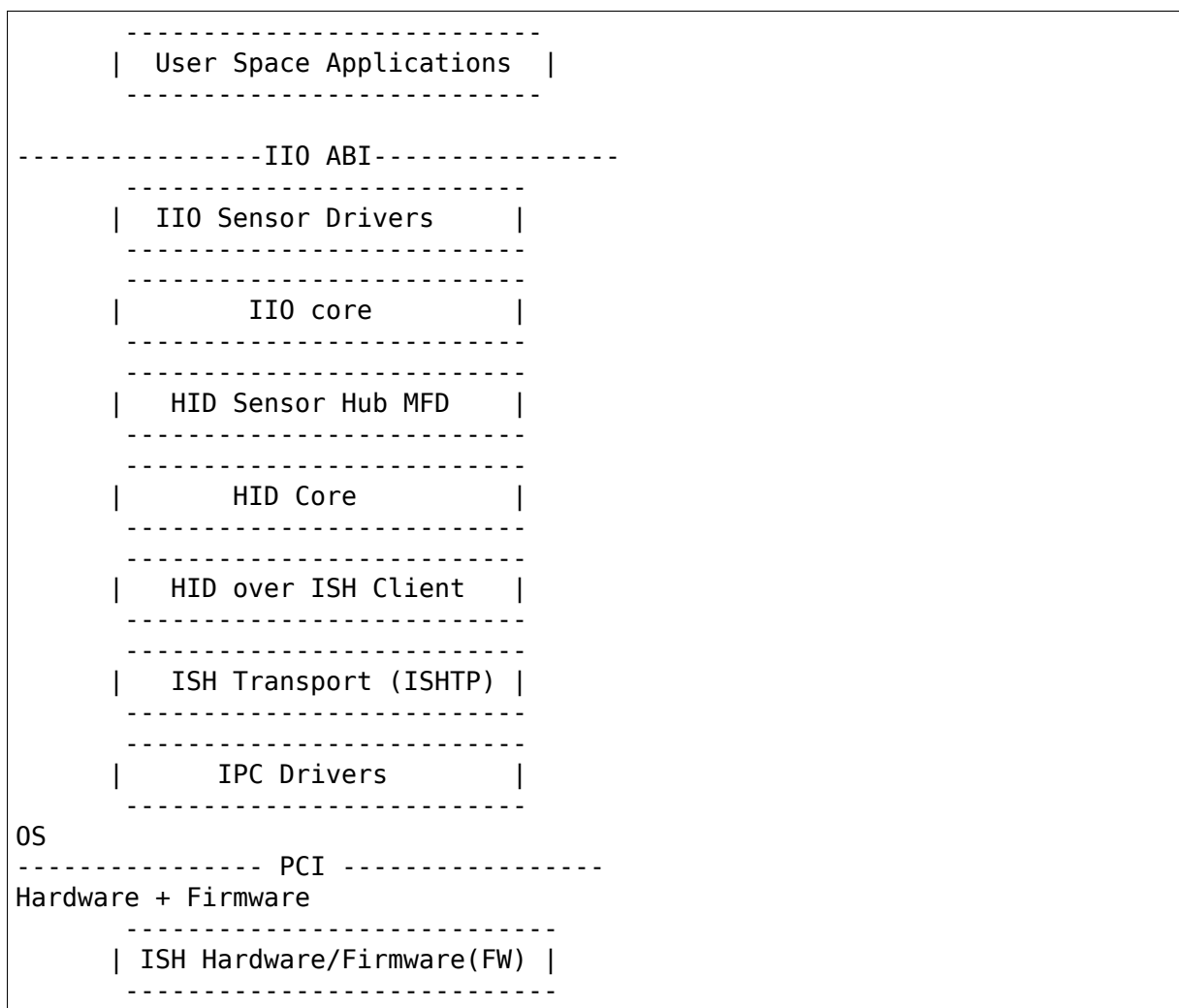
Like USB protocol provides a method for device enumeration, link management and user data encapsulation, the ISH also provides similar services. But it is very

light weight tailored to manage and communicate with ISH client applications implemented in the firmware.

The ISH allows multiple sensor management applications executing in the firmware. Like USB endpoints the messaging can be to/from a client. As part of enumeration process, these clients are identified. These clients can be simple HID sensor applications, sensor calibration application or sensor firmware update application.

The implementation model is similar, like USB bus, ISH transport is also implemented as a bus. Each client application executing in the ISH processor is registered as a device on this bus. The driver, which binds each device (ISH HID driver) identifies the device type and registers with the hid core.

7.2 2. ISH Implementation: Block Diagram



7.3 3. High level processing in above blocks

7.3.1 3.1 Hardware Interface

The ISH is exposed as “Non-VGA unclassified PCI device” to the host. The PCI product and vendor IDs are changed from different generations of processors. So the source code which enumerate drivers needs to update from generation to generation.

7.3.2 3.2 Inter Processor Communication (IPC) driver

Location: drivers/hid/intel-ish-hid/ipc

The IPC message used memory mapped I/O. The registers are defined in hw-ish-regs.h.

3.2.1 IPC/FW message types

There are two types of messages, one for management of link and other messages are to and from transport layers.

TX and RX of Transport messages

A set of memory mapped register offers support of multi byte messages TX and RX (E.g. IPC_REG_ISH2HOST_MSG, IPC_REG_HOST2ISH_MSG). The IPC layer maintains internal queues to sequence messages and send them in order to the FW. Optionally the caller can register handler to get notification of completion. A door bell mechanism is used in messaging to trigger processing in host and client firmware side. When ISH interrupt handler is called, the ISH2HOST doorbell register is used by host drivers to determine that the interrupt is for ISH.

Each side has 32 32-bit message registers and a 32-bit doorbell. Doorbell register has the following format: Bits 0..6: fragment length (7 bits are used) Bits 10..13: encapsulated protocol Bits 16..19: management command (for IPC management protocol) Bit 31: doorbell trigger (signal H/W interrupt to the other side) Other bits are reserved, should be 0.

3.2.2 Transport layer interface

To abstract HW level IPC communication, a set of callbacks are registered. The transport layer uses them to send and receive messages. Refer to struct `ishtp_hw_ops` for callbacks.

7.3.3 3.3 ISH Transport layer

Location: drivers/hid/intel-ish-hid/ishtp/

3.3.1 A Generic Transport Layer

The transport layer is a bi-directional protocol, which defines: - Set of commands to start, stop, connect, disconnect and flow control (ishtp/hbm.h) for details - A flow control mechanism to avoid buffer overflows

This protocol resembles bus messages described in the following document: <http://www.intel.com/content/dam/www/public/us/en/documents/technical-specifications/dcmi-hi-1-0-spec.pdf> “Chapter 7: Bus Message Layer”

3.3.2 Connection and Flow Control Mechanism

Each FW client and a protocol is identified by an UUID. In order to communicate to a FW client, a connection must be established using connect request and response bus messages. If successful, a pair (host_client_id and fw_client_id) will identify the connection.

Once connection is established, peers send each other flow control bus messages independently. Every peer may send a message only if it has received a flow-control credit before. Once it sent a message, it may not send another one before receiving the next flow control credit. Either side can send disconnect request bus message to end communication. Also the link will be dropped if major FW reset occurs.

3.3.3 Peer to Peer data transfer

Peer to Peer data transfer can happen with or without using DMA. Depending on the sensor bandwidth requirement DMA can be enabled by using module parameter ishtp_use_dma under intel_ishtp.

Each side (host and FW) manages its DMA transfer memory independently. When an ISHTP client from either host or FW side wants to send something, it decides whether to send over IPC or over DMA; for each transfer the decision is independent. The sending side sends DMA_XFER message when the message is in the respective host buffer (TX when host client sends, RX when FW client sends). The recipient of DMA message responds with DMA_XFER_ACK, indicating the sender that the memory region for that message may be reused.

DMA initialization is started with host sending DMA_ALLOC_NOTIFY bus message (that includes RX buffer) and FW responds with DMA_ALLOC_NOTIFY_ACK. Additionally to DMA address communication, this sequence checks capabilities: if the host doesn't support DMA, then it won't send DMA allocation, so FW can't send DMA; if FW doesn't support DMA then it won't respond with DMA_ALLOC_NOTIFY_ACK, in which case host will not use DMA transfers. Here ISH acts as busmaster DMA controller. Hence when host sends DMA_XFER, it's request to do host->ISH DMA transfer; when FW sends DMA_XFER, it means

that it already did DMA and the message resides at host. Thus, DMA_XFER and DMA_XFER_ACK act as ownership indicators.

At initial state all outgoing memory belongs to the sender (TX to host, RX to FW), DMA_XFER transfers ownership on the region that contains ISHTP message to the receiving side, DMA_XFER_ACK returns ownership to the sender. A sender needs not wait for previous DMA_XFER to be ack'ed, and may send another message as long as remaining continuous memory in its ownership is enough. In principle, multiple DMA_XFER and DMA_XFER_ACK messages may be sent at once (up to IPC MTU), thus allowing for interrupt throttling. Currently, ISH FW decides to send over DMA if ISHTP message is more than 3 IPC fragments and via IPC otherwise.

3.3.4 Ring Buffers

When a client initiate a connection, a ring or RX and TX buffers are allocated. The size of ring can be specified by the client. HID client set 16 and 32 for TX and RX buffers respectively. On send request from client, the data to be sent is copied to one of the send ring buffer and scheduled to be sent using bus message protocol. These buffers are required because the FW may have not have processed the last message and may not have enough flow control credits to send. Same thing holds true on receive side and flow control is required.

3.3.5 Host Enumeration

The host enumeration bus command allow discovery of clients present in the FW. There can be multiple sensor clients and clients for calibration function.

To ease in implantation and allow independent driver handle each client this transport layer takes advantage of Linux Bus driver model. Each client is registered as device on the the transport bus (ishtp bus).

Enumeration sequence of messages:

- Host sends HOST_START_REQ_CMD, indicating that host ISHTP layer is up.
- FW responds with HOST_START_RES_CMD
- Host sends HOST_ENUM_REQ_CMD (enumerate FW clients)
- FW responds with HOST_ENUM_RES_CMD that includes bitmap of available FW client IDs
- For each FW ID found in that bitmap host sends HOST_CLIENT_PROPERTIES_REQ_CMD
- FW responds with HOST_CLIENT_PROPERTIES_RES_CMD. Properties include UUID, max ISHTP message size, etc.
- Once host received properties for that last discovered client, it considers ISHTP device fully functional (and allocates DMA buffers)

7.3.4 3.4 HID over ISH Client

Location: drivers/hid/intel-ish-hid

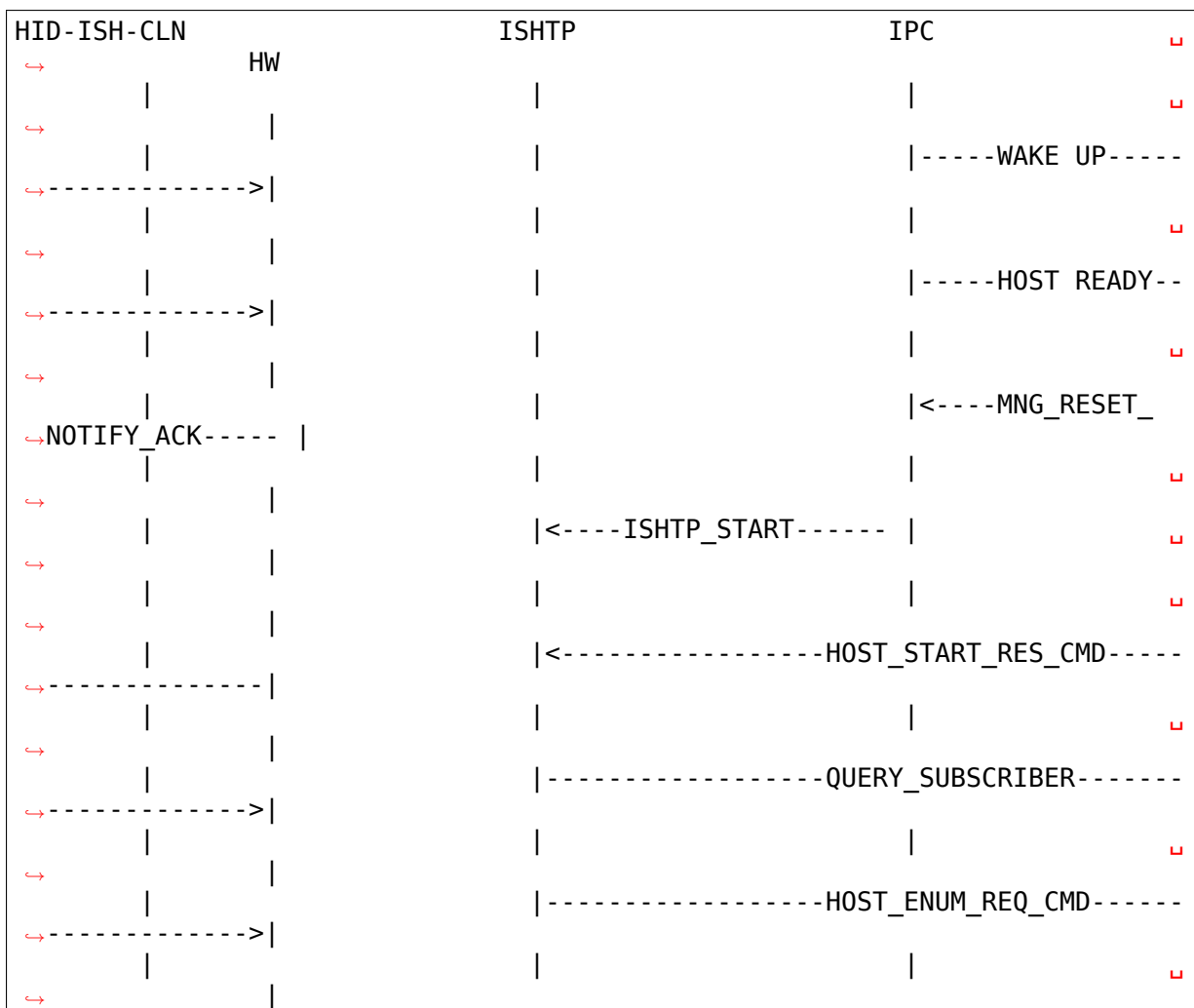
The ISHTP client driver is responsible for:

- enumerate HID devices under FW ISH client
- Get Report descriptor
- Register with HID core as a LL driver
- Process Get/Set feature request
- Get input reports

7.3.5 3.5 HID Sensor Hub MFD and IIO sensor drivers

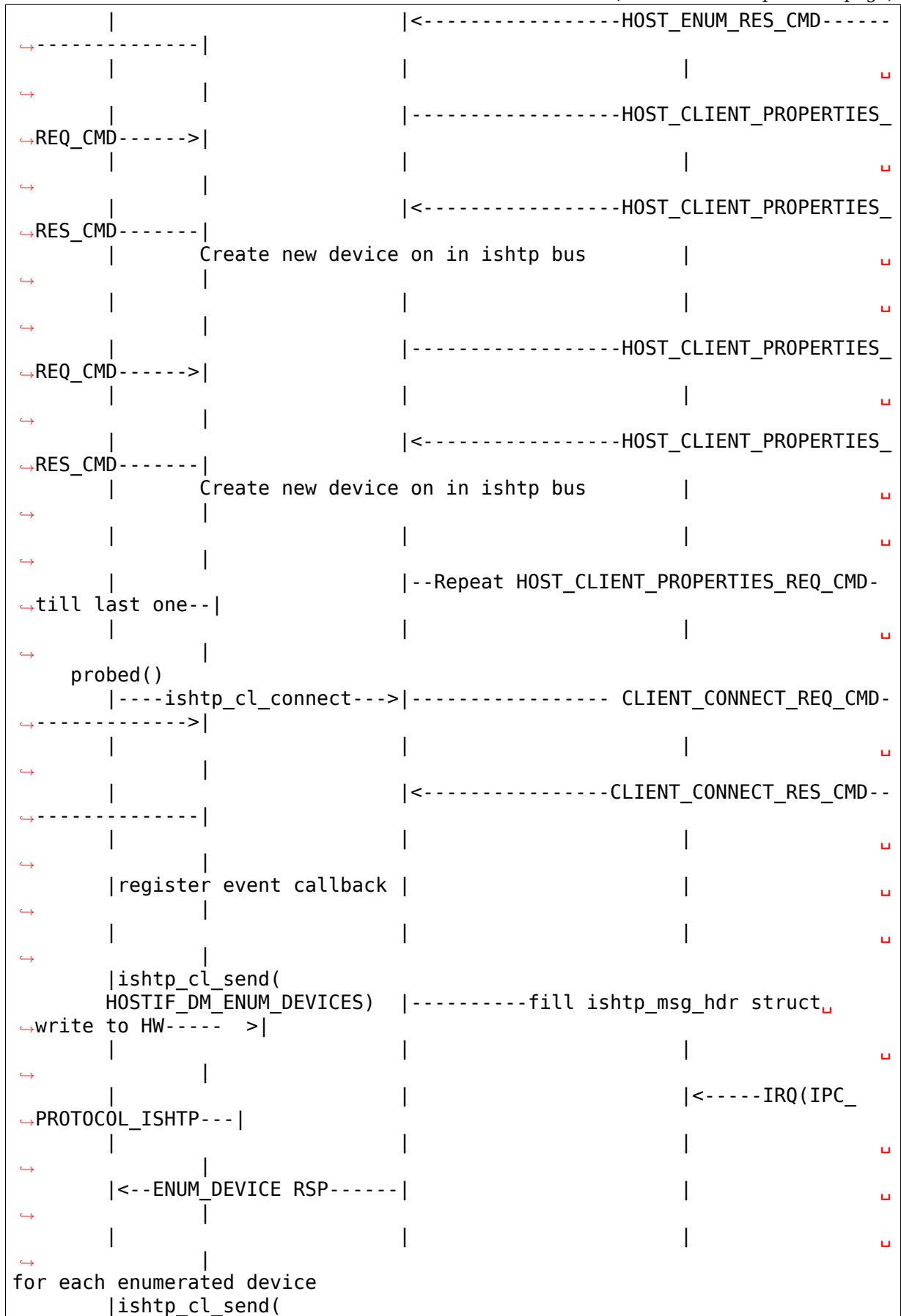
The functionality in these drivers is the same as an external sensor hub. Refer to Documentation/hid/hid-sensor.rst for HID sensor Documentation/ABI/testing/sysfs-bus-iio for IIO ABIs to user space

7.3.6 3.6 End to End HID transport Sequence Diagram



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```

HOSTIF_GET_HID_DESCRIPTOR|-----fill ishtp_msg_hdr struct
↪write to HW----- >|
|
|
↪
|
|
...Response
|
|
↪
for each enumerated device
|ishtp_cl_send(
HOSTIF_GET_REPORT_DESCRIPTOR|-----fill ishtp_msg_hdr struct
↪write to HW-- >|
|
|
↪
|
|
hid_allocate_device
|
|
↪
hid_add_device
|
|
↪
|
|
↪

```

7.3.7 3.7 ISH Debugging

To debug ISH, event tracing mechanism is used. To enable debug logs echo 1 > /sys/kernel/debug/tracing/events/intel_ish/enable cat /sys/kernel/debug/tracing/trace

7.3.8 3.8 ISH IIO sysfs Example on Lenovo thinkpad Yoga 260

```

root@otcpl-ThinkPad-Yoga-260:~# tree -l /sys/bus/iio/devices/
/sys/bus/iio/devices/
├── iio:device0 -> ../../../../devices/0044:8086:22D8.0001/HID-SENSOR-200073.
↪9.auto/iio:device0
|
| ├── buffer
| | ├── enable
| | ├── length
| | └── watermark
|
| ...
| ├── in_accel_hysteresis
| ├── in_accel_offset
| ├── in_accel_sampling_frequency
| ├── in_accel_scale
| ├── in_accel_x_raw
| ├── in_accel_y_raw
| ├── in_accel_z_raw
| ├── name
| ├── scan_elements
| | ├── in_accel_x_en
| | └── in_accel_x_index
|

```

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	<ul style="list-style-type: none"> └─ in_accel_x_type └─ in_accel_y_en └─ in_accel_y_index └─ in_accel_y_type └─ in_accel_z_en └─ in_accel_z_index └─ in_accel_z_type
...	<ul style="list-style-type: none"> └─ devices <ul style="list-style-type: none"> └─ buffer <ul style="list-style-type: none"> └─ enable └─ length └─ watermark └─ dev └─ in_intensity_both_raw └─ in_intensity_hysteresis └─ in_intensity_offset └─ in_intensity_sampling_frequency └─ in_intensity_scale └─ name └─ scan_elements <ul style="list-style-type: none"> └─ in_intensity_both_en └─ in_intensity_both_index └─ in_intensity_both_type └─ trigger <ul style="list-style-type: none"> └─ current_trigger
...	<ul style="list-style-type: none"> └─ buffer <ul style="list-style-type: none"> └─ enable └─ length └─ watermark └─ dev └─ in_magn_hysteresis └─ in_magn_offset └─ in_magn_sampling_frequency └─ in_magn_scale └─ in_magn_x_raw └─ in_magn_y_raw └─ in_magn_z_raw └─ in_rot_from_north_magnetic_tilt_comp_raw └─ in_rot_hysteresis └─ in_rot_offset └─ in_rot_sampling_frequency └─ in_rot_scale └─ name
...	<ul style="list-style-type: none"> └─ scan_elements <ul style="list-style-type: none"> └─ in_magn_x_en └─ in_magn_x_index └─ in_magn_x_type └─ in_magn_y_en └─ in_magn_y_index

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			<ul style="list-style-type: none"> └─ in_magn_y_type └─ in_magn_z_en └─ in_magn_z_index └─ in_magn_z_type └─ in_rot_from_north_magnetic_tilt_comp_en └─ in_rot_from_north_magnetic_tilt_comp_index └─ in_rot_from_north_magnetic_tilt_comp_type └─ trigger <ul style="list-style-type: none"> └─ current_trigger
...			<ul style="list-style-type: none"> └─ buffer <ul style="list-style-type: none"> └─ enable └─ length └─ watermark └─ dev └─ in_anglvel_hysteresis └─ in_anglvel_offset └─ in_anglvel_sampling_frequency └─ in_anglvel_scale └─ in_anglvel_x_raw └─ in_anglvel_y_raw └─ in_anglvel_z_raw └─ name └─ scan_elements <ul style="list-style-type: none"> └─ in_anglvel_x_en └─ in_anglvel_x_index └─ in_anglvel_x_type └─ in_anglvel_y_en └─ in_anglvel_y_index └─ in_anglvel_y_type └─ in_anglvel_z_en └─ in_anglvel_z_index └─ in_anglvel_z_type └─ trigger <ul style="list-style-type: none"> └─ current_trigger
...			<ul style="list-style-type: none"> └─ buffer <ul style="list-style-type: none"> └─ enable └─ length └─ watermark └─ dev └─ in_anglvel_hysteresis └─ in_anglvel_offset └─ in_anglvel_sampling_frequency └─ in_anglvel_scale └─ in_anglvel_x_raw └─ in_anglvel_y_raw └─ in_anglvel_z_raw └─ name └─ scan_elements <ul style="list-style-type: none"> └─ in_anglvel_x_en └─ in_anglvel_x_index

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